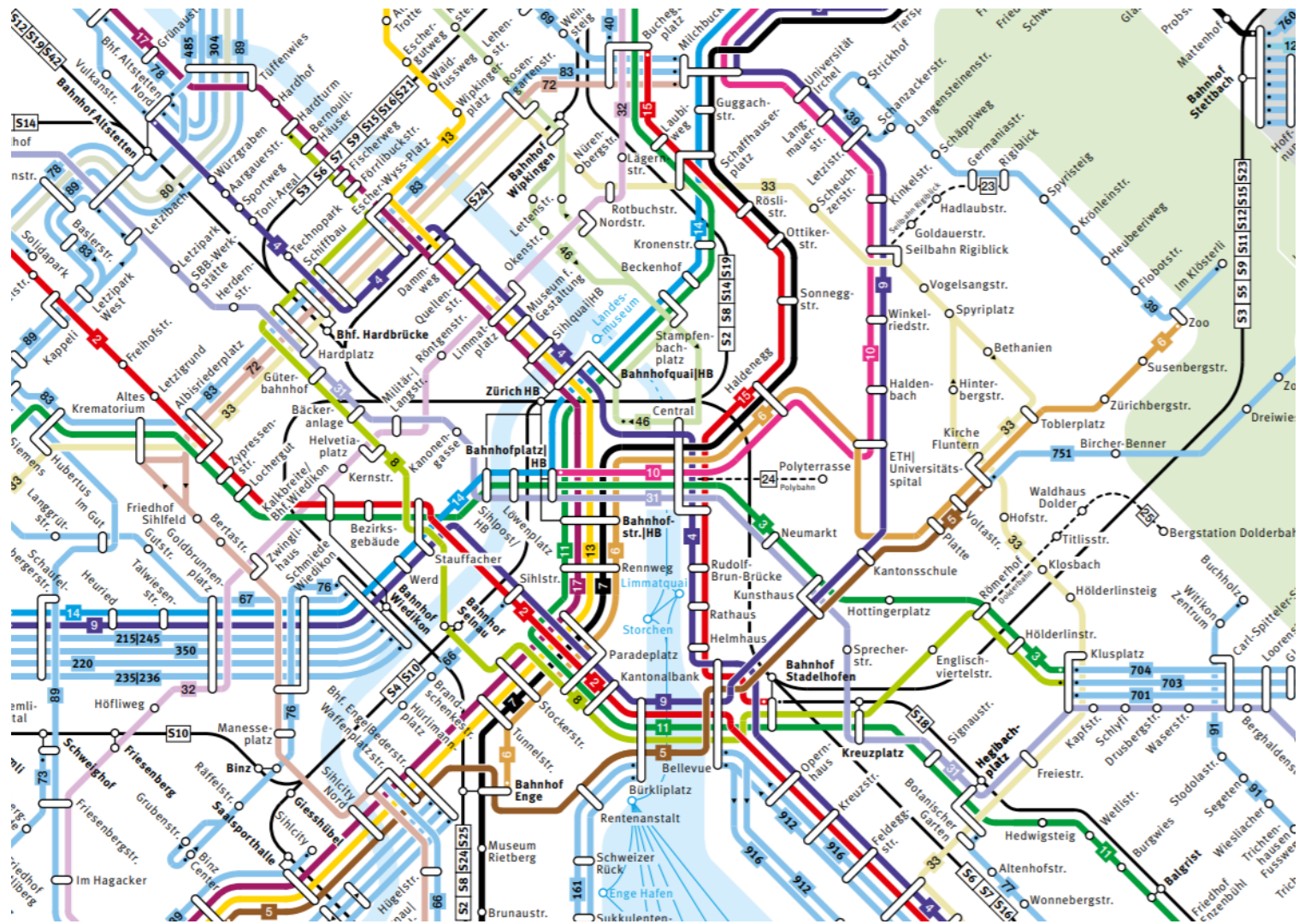


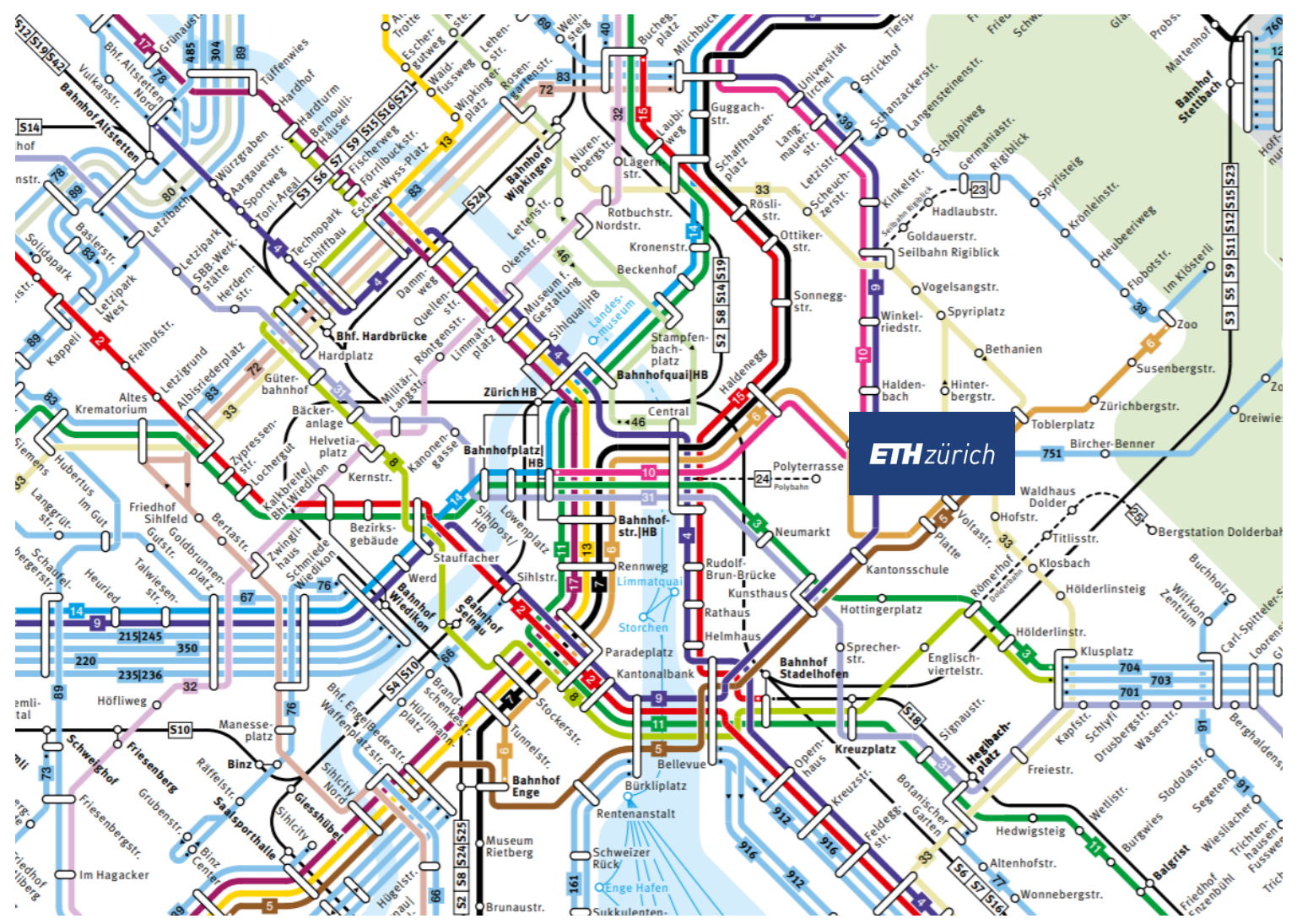
Congestion and Stretch Aware Static Fast Rerouting [appeared @INFOCOM'19]

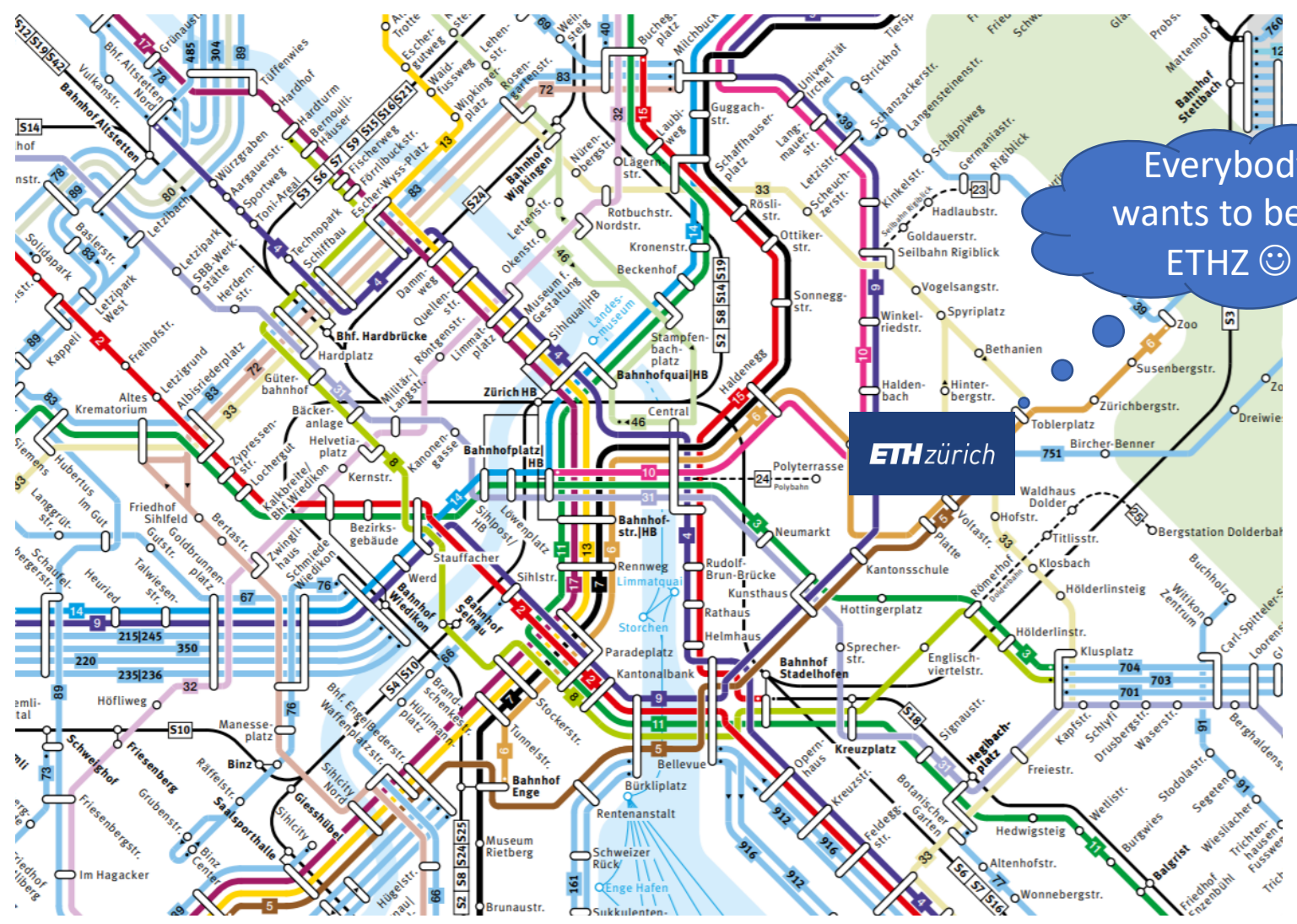
Klaus-Tycho Foerster, Yvonne-Anne Pignolet (DFINITY), Stefan Schmid, and Gilles Tredan (LAAS-CNRS)





[Idea taken from Gilles Tredan]

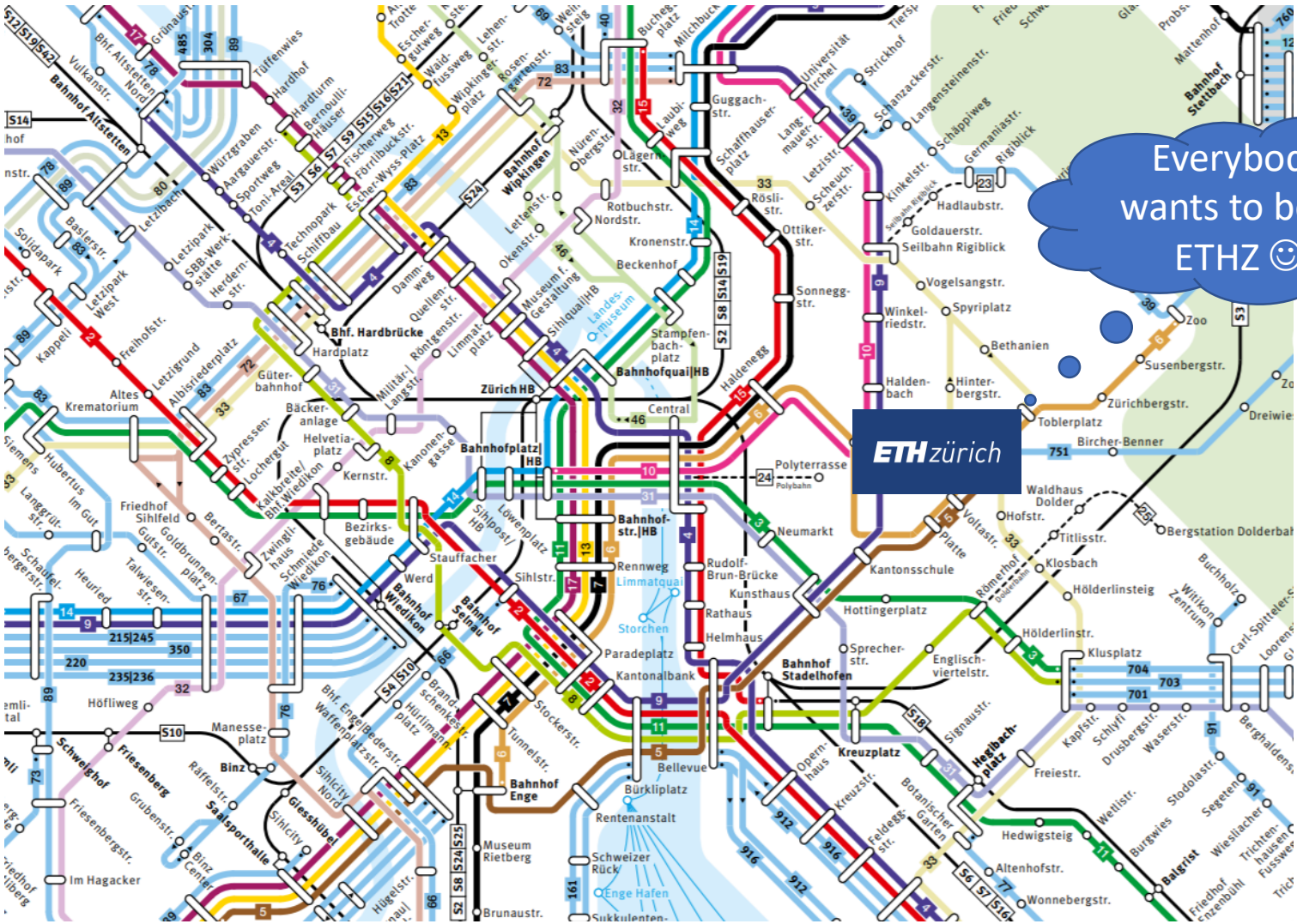




Everybody
wants to be at
ETHZ 😊

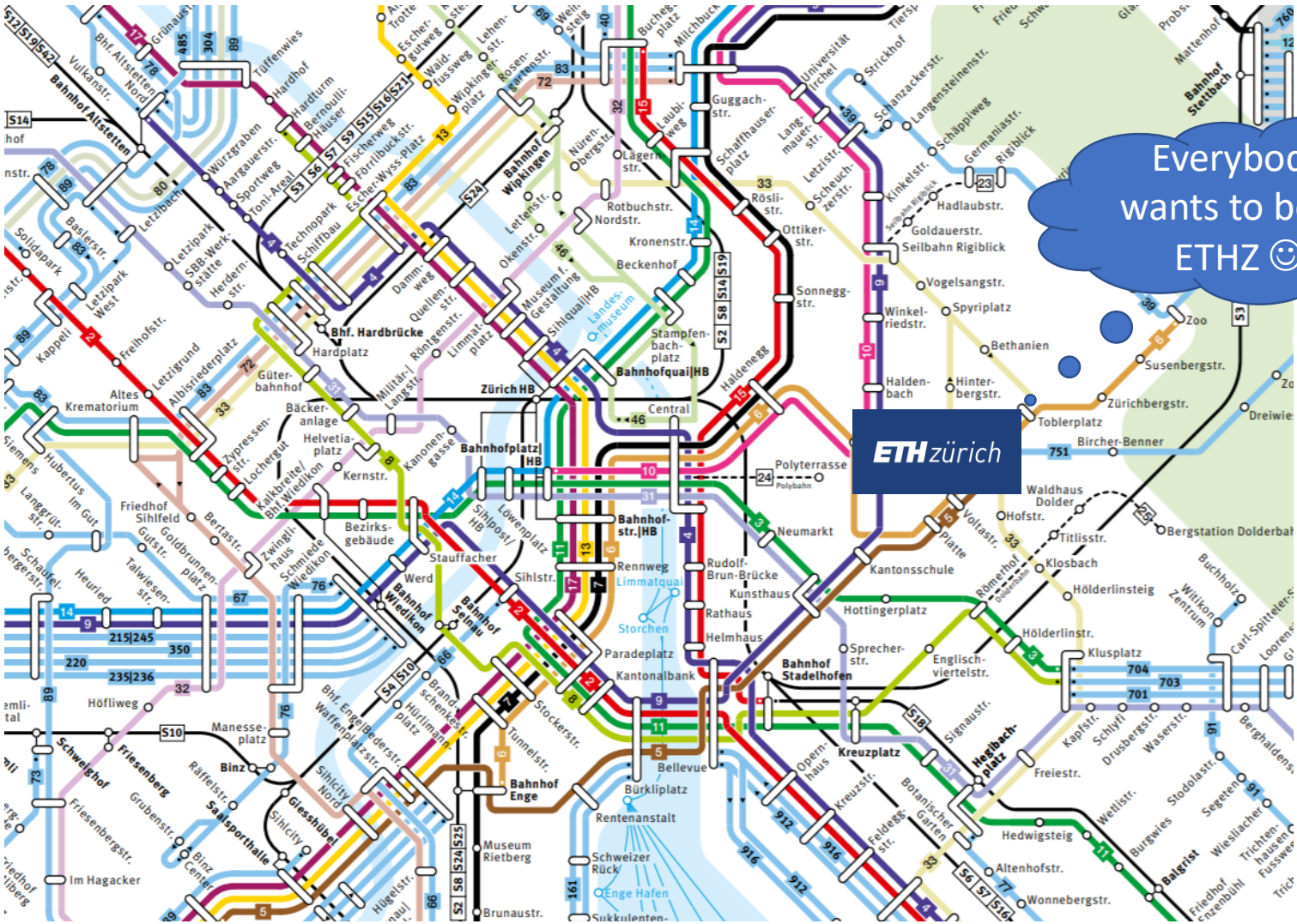
ETH zürich

What if a link fails?



Everybody wants to be at ETHZ 😊

What if a link fails? Take a detour 😊



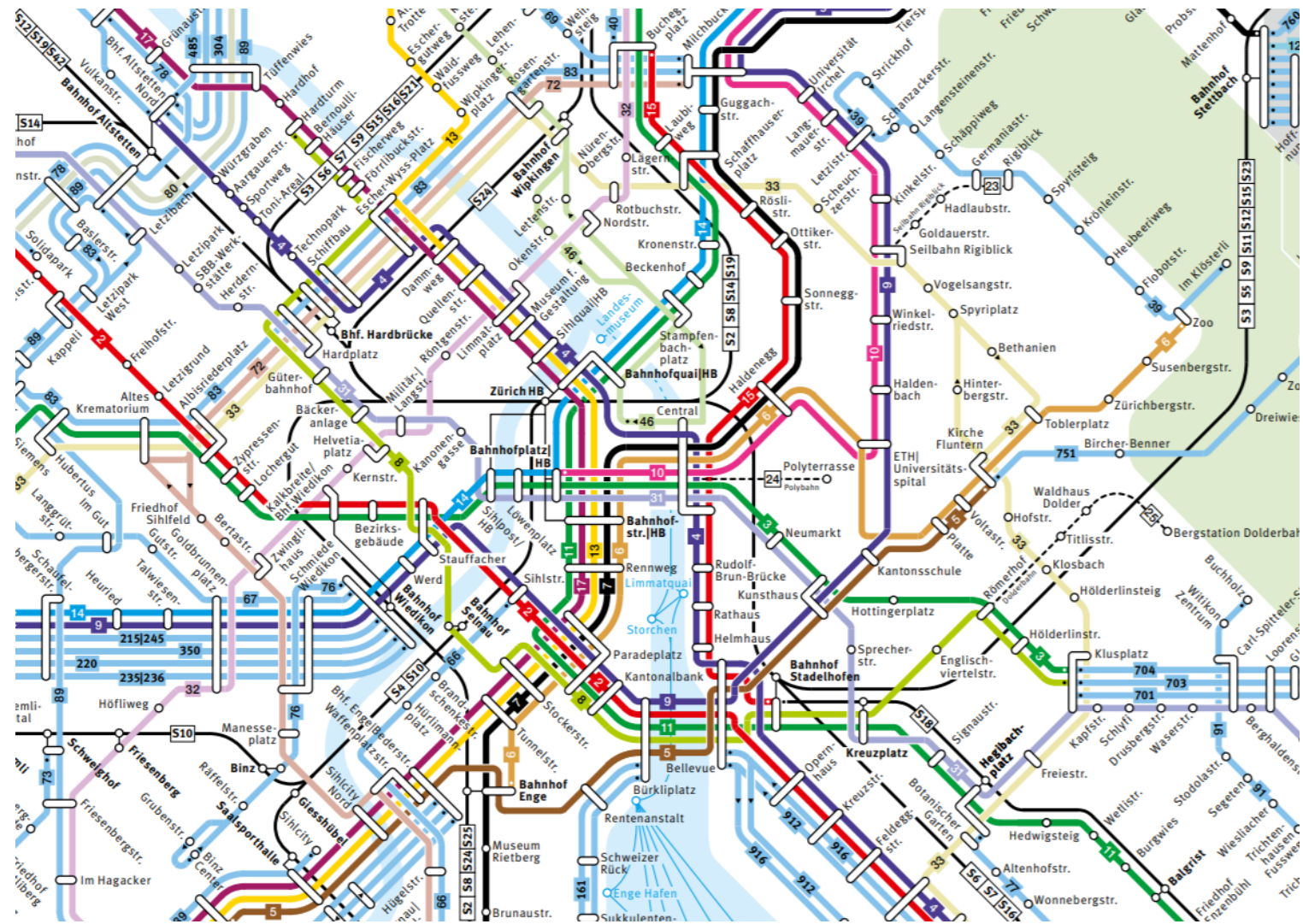
Everybody wants to be at ETHZ 😊

Everybody takes the same detour? High load!



Distribute people over all detours? High path stretch!



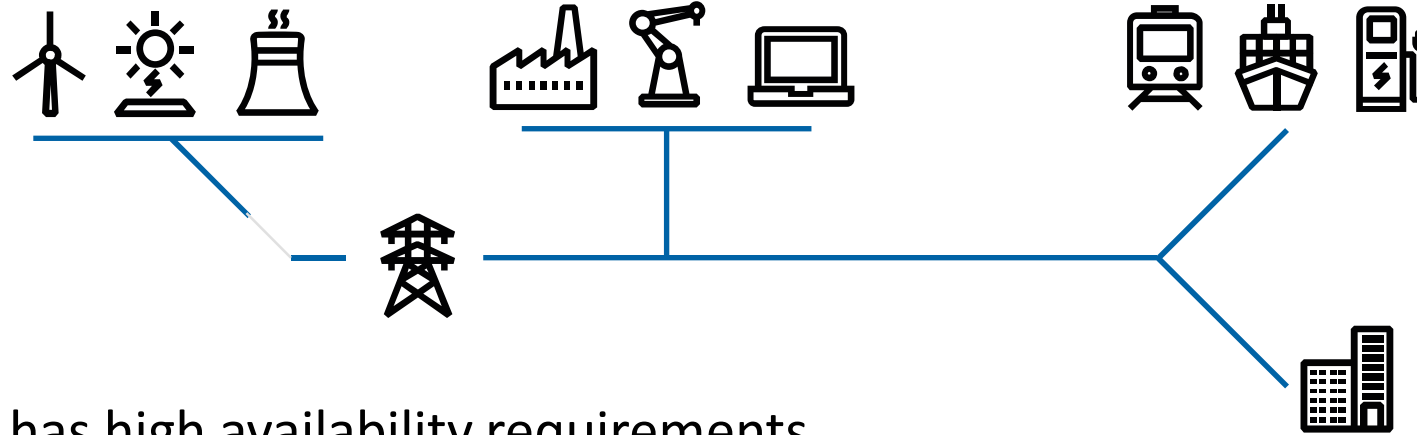




"The disparity in timescales between packet forwarding (which can be less than a microsecond) and control plane convergence (which can be as high as hundreds of milliseconds) means that failures often lead to unacceptably long outages"

Ensuring Connectivity via Data Plane Mechanisms: NSDI'13

Motivation



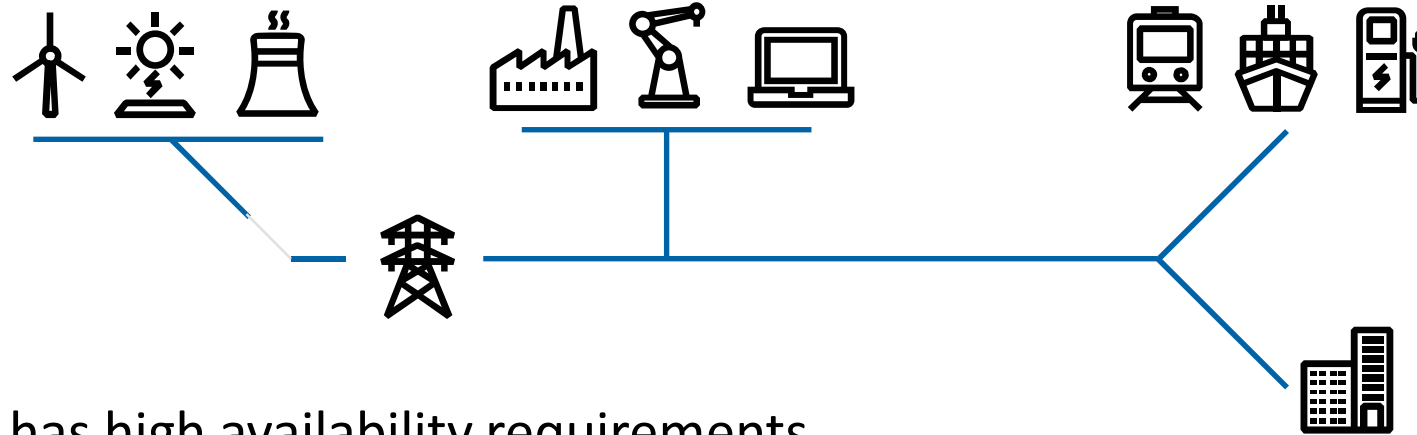
- Critical infrastructure has high availability requirements
- Industrial systems are more and more connected
- Hard real-time requirements



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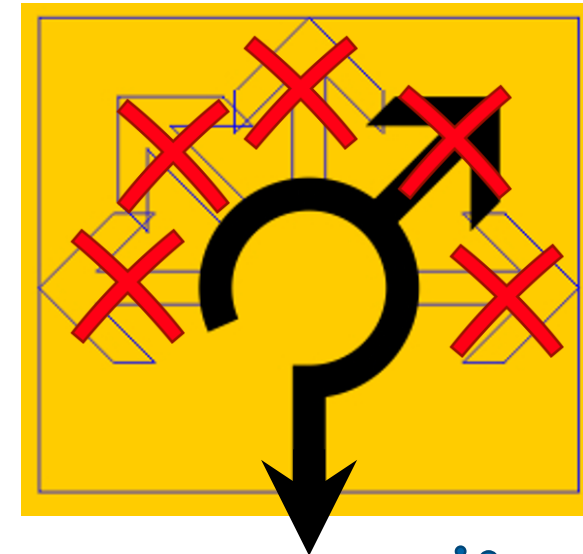
⇒ How to provide dependability guarantee despite link failures in networks?
⇒ Possible without communication between nodes?
⇒ With low load? With low stretch?

Talk Structure

1. Model and Objectives
2. Background and Lower Bounds
3. Algorithms and Upper Bounds
4. Simulation Results
5. Conclusion and Outlook

Model I/II: Routing and Network

- Network is a strongly connected directed graph
- Forwarding may only match on:
 1. Source
 2. Destination
 3. Incident failures
 4. Incoming port
- No packet (header) changes allowed, no communication
- Static routing tables, deterministic behaviour
- Single destination routing, uniform flow sizes



Route can be
a *walk*

Model II/II: Quality *from a Worst-Case Perspective*

1. Resilience

- How many link failures can we survive and still guarantee delivery?
- Upper bound: $(r+1)$ -link-connected graph: at most r

2. Load

- Maximum additional link utilization due to rerouting

3. Stretch

- Maximum additional hops due to rerouting

The Price of Locality (for *every* Scheme and Graph)

Stretch under r failures:

- Adversary can force to visit $r+1$ neighbors of destination •

Fail r links incident to the destination

Load under r failures:

- Adversary can force additional load of \sqrt{r} • • •

Previously only weaker bound known,
without incoming port

Let's try to meet this bound for many flows

CASA: Rerouting on Arborescences

- Takes arborescences as input *e.g.* generated by Chiesa *et al.*
 - Influences the stretch, we get good bounds for *e.g.* so-called *independent spanning trees*

Algorithm

- 1: Determine current arborescence T from in-port
- 2: If next hop in T alive, use it, else
- 3: Pick next arborescence T' from **BIBD-Matrix**

until the next
hop is alive

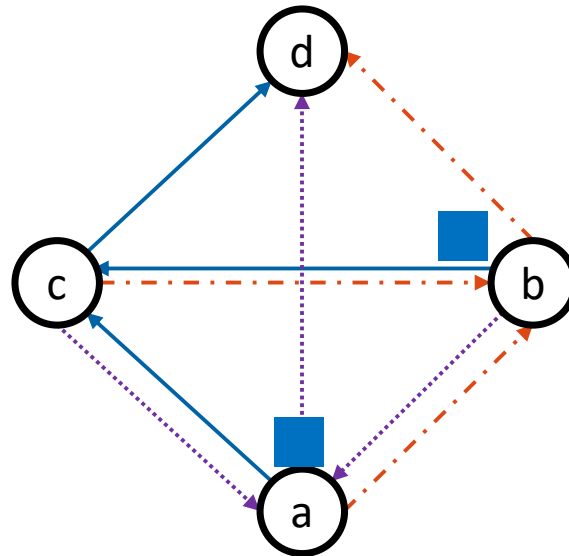
different flows
use different T'

1	4	12	9	3		5
3	12	4	5	6	...	8
4	6	13	8	11		
x	x	x		4		
x	x	x	x	x	4	9

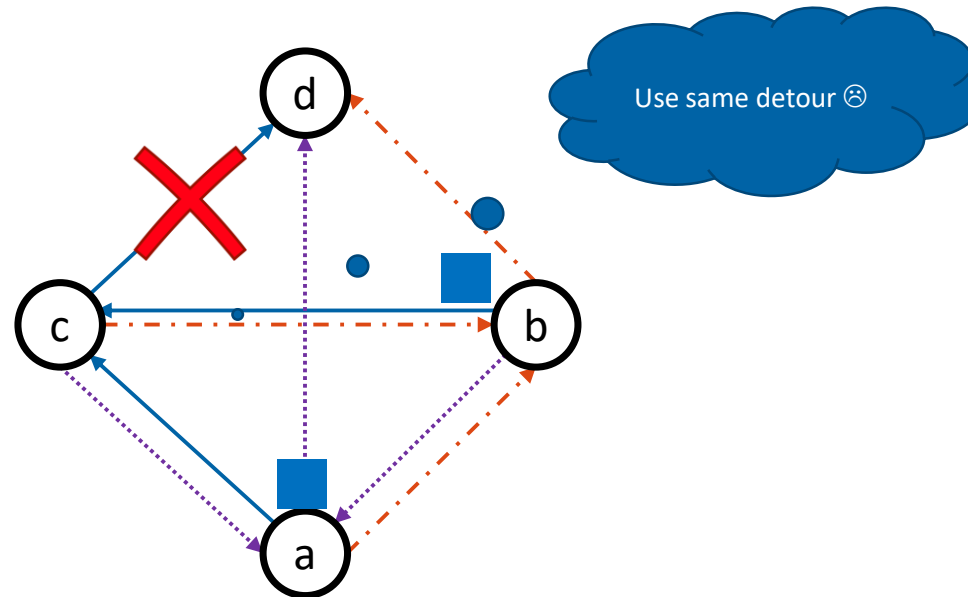


We re-structure BIBD-matrix to be good for many flows

CASA: Example *without* BIBD



CASA: Example *without* BIBD

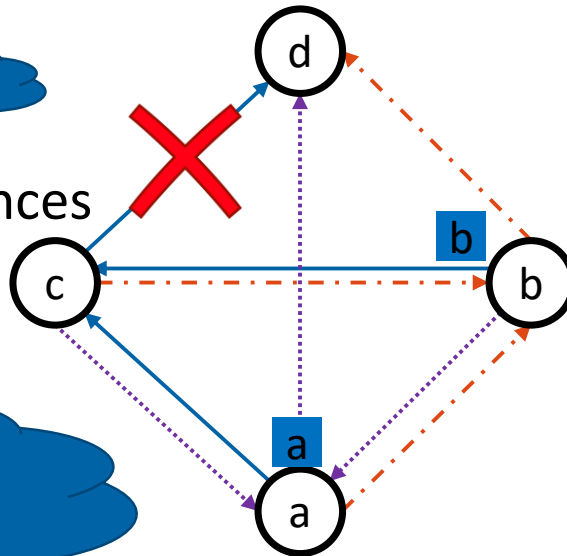


CASA: Example *with* BIBD

How much extra load?

- Up to $O(\sqrt{r})$. . .
- For more flows than #arborescences

Lower bound: \sqrt{r}



$$\sqrt{\#failures} < \frac{(\#arborescences)^{\frac{3}{2}}}{\#flows}$$

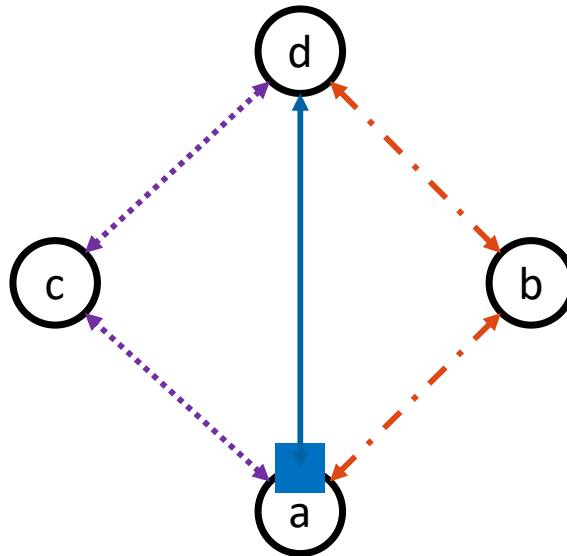
Beyond CASA

- $r+1$ arborescences give r -resiliency under directed link failures
 - But unclear how to obtain r -resiliency under bi-directed link failures
- Motivation for a simplified heuristic: *SquareOne*
 - Pick $r+1$ bi-directed link-disjoint source-destination paths
 - Under failure: bounce back to the source, pick next path



<https://Netflix.com>

SquareOne



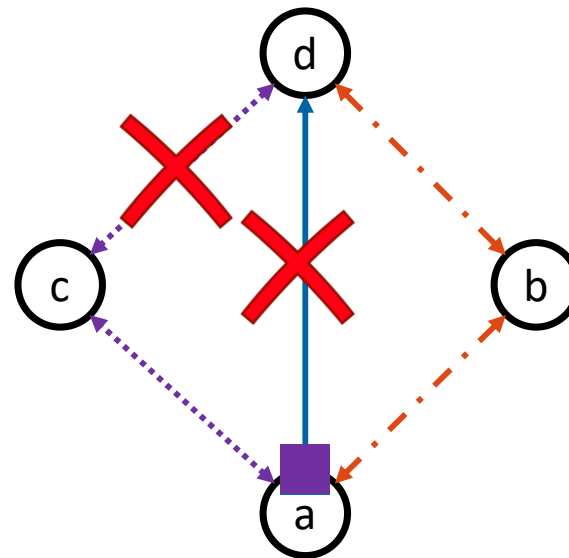
SquareOne



<https://Netflix.com>

How good in practice?

No theoretical guarantees beyond resiliency



Easy to compute via *e.g.* max-flow formulations. Order path priority *e.g.* by length

Selected Evaluations

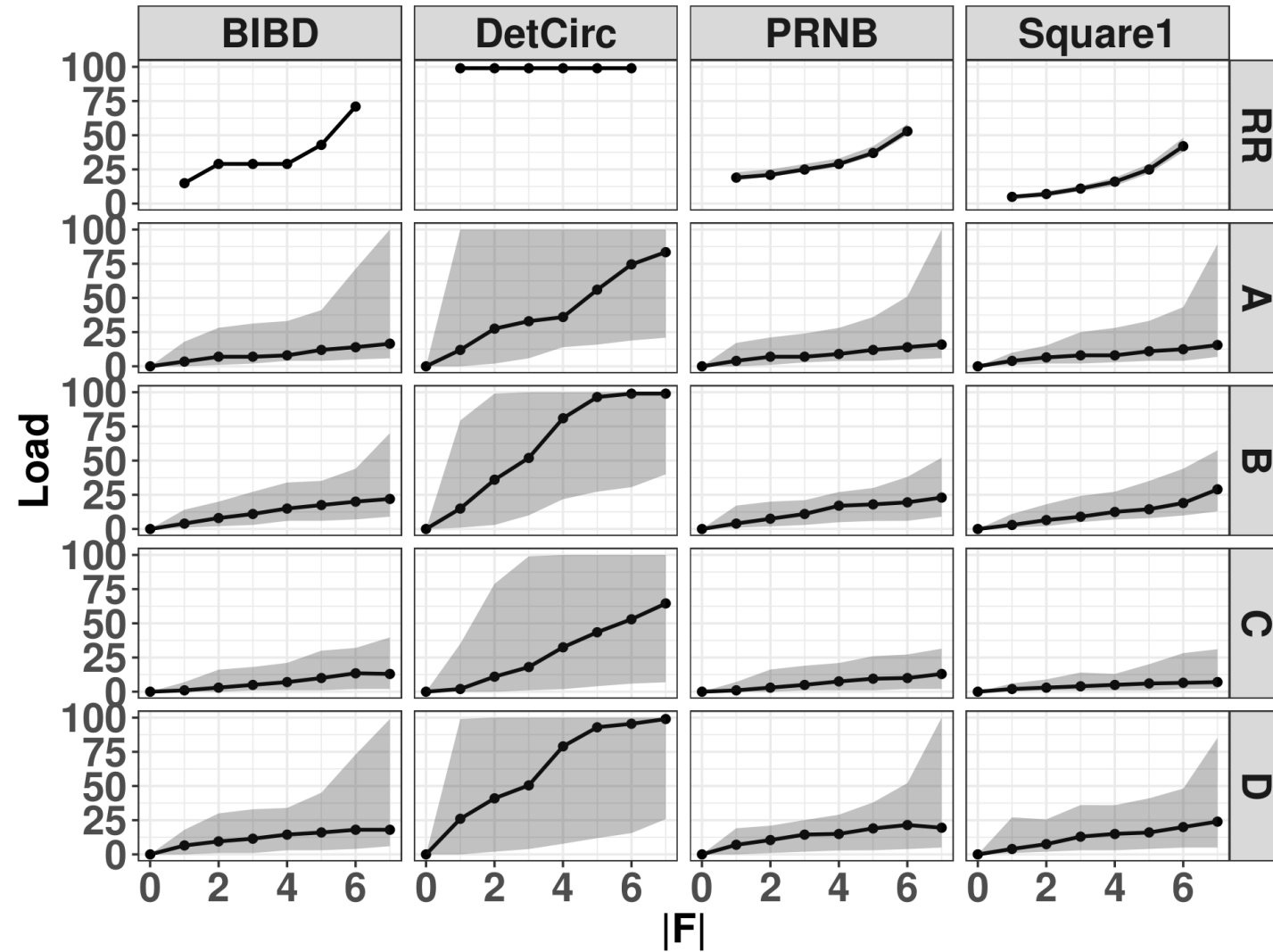
- 8-connected 8-regular random graphs (**RR**, 100 routers each)
- well-connected cores of real-world ASes (*Rocketfuel*) (204-387 routers, 1667-4736 links)
- Three arborescence methods (using the *same* arborescences)
 - **CASA (BIBD)**
 - Deterministic Circular (**DetCirc**) from Chiesa *et al.*
 - Random (**PRNB**) from Chiesa *et al.*
- Also: **SquareOne**

Setting from prior work

Thanks to Marco Chiesa and Ilya Nikolaevskiy for their support

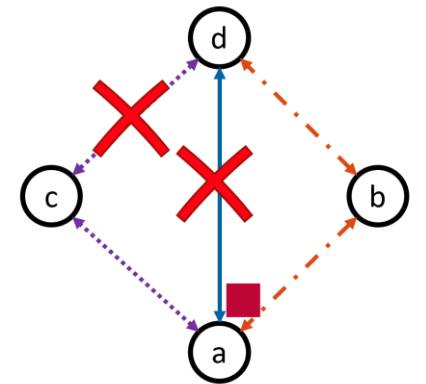
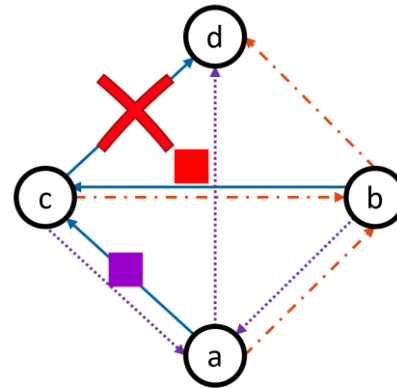
Issues in practice:
Real randomness on routers?
Packet reordering?

Deterministic Worst-Case Failures



Conclusion

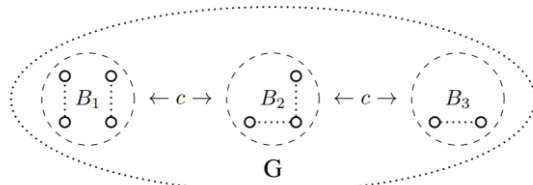
- We present **efficient static fast failover schemes** on general graphs
 - **CASA**: Combines **arborescences** and improved **block-designs** (BIBDs)
 - With theoretical guarantees
 - **SquareOne**: Well performing resilient heuristic
 - Based on edge-disjoint paths
- Next slide: Further related problems we work on



Some More Related Problems

- Improving arborescence decompositions

- #1: Build small stretch arborescences in parallel
 - Current approach: build sequentially in greedy fashion
 - Benefit: Resilient to more failures under nice distributions



- #2: Account for e.g. Shared Risk Link Groups (SRLGs)
 - Leverage post-processing according to objective function
 - Ideally: A SRLG is contained in a single arborescence

Appears at #1: DSN 2019, #2: SRDS 2019

- Allowing packet header modification (MPLS, SR)

- #1: More powerful, but harder to verify correctness?
 - MPLS w. multiple link failures: verification in polynomial time!

	P-Rex	NetKAT	HSA	VeriFlow	Anteater
Protocol Support	SR/MPLS	OF	Agn.	OF	Agn.
Approach	Autom.	Alg.	Geom.	Tries	SAT
Complexity	Polynom.	PSPACE	Polynom.	NP	NP
Static	✓	✓	✓	✗	✓
Reachability	✓	✓	✓	✓	✓
Loop Queries	✓	✓	✓	✓	✓
What-if	✓	N/A	✓	N/A	✗
Unlim. Header	✓	N/A	✗	✗	N/A
Performance	✓	✓ [1]	✓	✓	✓
Waypointing	✓	✓	✓	✓	✗
Language	Py., C	OCaml	Py., C	Py.	C++, Ruby

- #2: Leverage Segment Routing (in Linux kernel for IPv6)
 - Allows maximal link protection e.g. in Hypercubes

Appears at #1: CoNEXT 2018, #2: OPODIS 2018

Papers

- *Improved Fast Rerouting Using Postprocessing*
Klaus-T. Foerster, Andrzej Kamisinski, Yvonne-Anne Pignolet, Stefan Schmid, and Gilles Tredan. *SRDS 2019*
- *Bonsai: Efficient Fast Failover Routing Using Small Arborescences*
Klaus-T. Foerster, Andrzej Kamisinski, Yvonne-Anne Pignolet, Stefan Schmid, and Gilles Tredan. *DSN 2019*
- *CASA: Congestion and Stretch Aware Static Fast Rerouting*
Klaus-T. Foerster, Yvonne-Anne Pignolet, Stefan Schmid, and Gilles Tredan. *INFOCOM 2019*
- *P-Rex: Fast Verification of MPLS Networks with Multiple Link Failures*
Jesper S. Jensen, Troels B. Krogh, Jonas S. Madsen, S. Schmid, Jiri Srba, and Marc T. Thorngersen. *CoNEXT 2018*
- *Local Fast Segment Rerouting on Hypercubes*
Klaus-T. Foerster, Mahmoud Parham, Stefan Schmid, and Tao Wen. *OPODIS 2018*

Congestion and Stretch Aware Static Fast Rerouting [appeared @INFOCOM'19]

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Papers Referenced

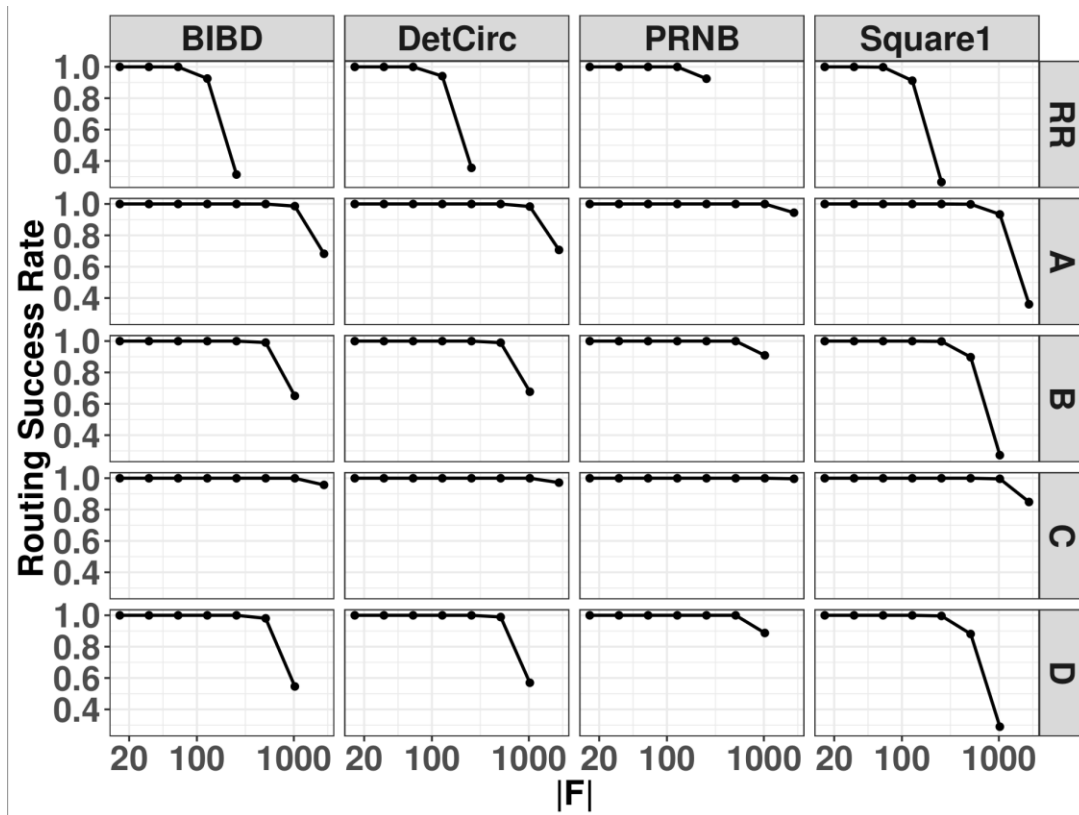
- *How (Not) to Shoot in Your Foot with SDN Local Fast Failover: A Load-Connectivity Tradeoff*
Michael Borokhovich and Stefan Schmid. *OPODIS 2013*
- *Load-Optimal Local Fast Rerouting for Dependable Networks*
Yvonne-Anne Pignolet, Stefan Schmid, and Gilles Tredan. *DSN 2013*
- *IP Fast Rerouting for Multi-Link Failures*
Theodore Elhourani, Abishek Gopalan, Srinivasan Ramasubramanian.
IEEE/ACM Trans. Netw. 24(5): 3014-3025 (2016)
- *The Quest for Resilient (Static) Forwarding Tables*
Marco Chiesa and Ilya Nikolaevskiy et al. *INFOCOM 2016*

Rocketfuel ASes

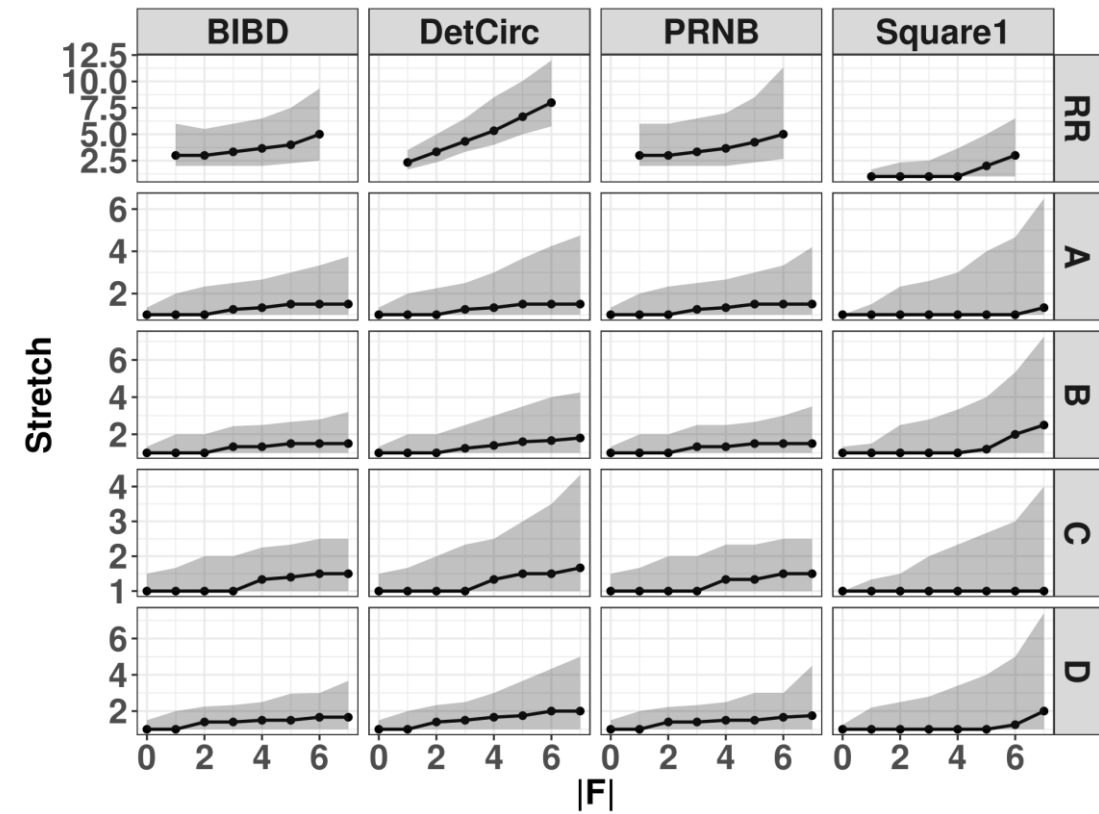
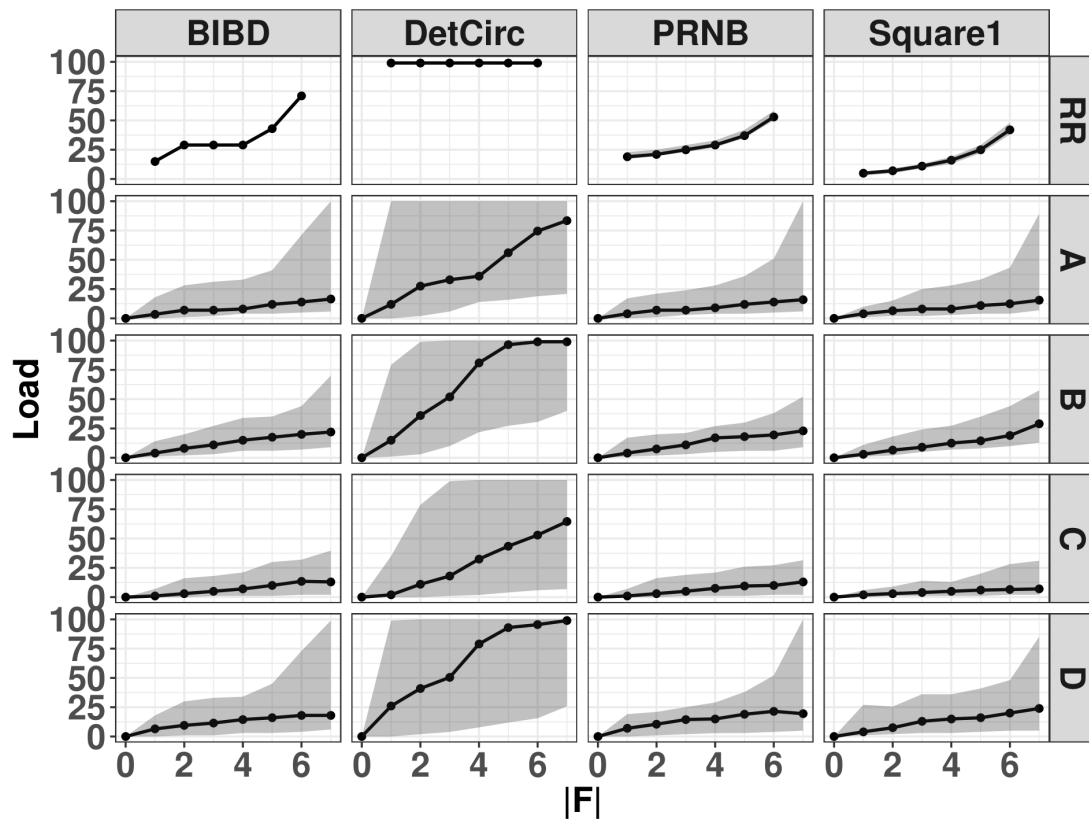
AS	1239 A	2914 B	3356 C	7018 D
Number of nodes	389	225	377	204
Number of links	3621	1696	4736	1667
Eccentricity	6	6	6	6
Avg shortest path length	3.06	2.48	3.14	3.17

TABLE I: Properties of 8-connected cores of various ASes

Evaluation: Resiliency



Evaluation: Deterministic Worst-Case Failures



Evaluation: Random Failures

